

## Claims

What we claim is:

1. In a shared memory model system, a method whereby, in a state wherein a plurality of threads exist, a bit that represents a lock type and an identifier for a thread that has acquired a lock in accordance with a first lock type, or an identifier of a second lock type, are stored in a storage area that corresponds to an object and a lock on an object is thus managed, said method comprising:

determining, if a second thread attempts to acquire a lock on a specific object that is held by a first thread, whether a bit that represents said lock type on said specific object represents said first lock type;

setting a contention bit if said bit represents said first lock type;

determining, before said first thread unlocks said specific object, whether said bit that represents said lock type represents said first lock type;

storing in said storage area a special identifier that differs from the identifiers for said plurality of threads;

issuing a synchronization command for said memory system;

storing in said storage area data indicating the absence of a thread that holds said lock on said specific object;

determining whether said contention bit has been set if said bit that represents said lock type represents said first lock type; and

terminating an unlocking process if said contention bit has not been set without any other process being performed.

2. The lock management method according to claim 1,  
further comprising:

shifting said first thread, when said contention  
bit has been set, to an exclusive control state for a  
mechanism that enables the exclusive control of the  
accessing of said object, and a thread waiting  
operation and the transmission to a waiting thread of  
a notification, both of which are to be performed  
when a predetermined condition has been established;

permitting said first thread to transmit said  
notification to said waiting thread;

setting said second thread in the busy waiting state,  
when said predetermined condition has not been  
established and when said special identifier has been  
stored, until a thread that holds said lock on said  
specific object is no longer present and until said  
bit that represents said lock type represents said  
first lock type; and

permitting said first thread to exit said  
exclusive control state.

3. The lock management method according to claim 1,  
wherein said first lock type is a lock method whereby to  
manage a lock state an identifier for a thread that has  
locked an object is stored in correlation with said  
object.



5. In a shared memory model system, an apparatus where, in a state wherein a plurality of threads exist, a bit that represents a lock type and an identifier for a thread that has acquired a lock in accordance with a first lock type, or an identifier of a second lock type, are stored in a storage area that corresponds to an object and a lock on an object is thus managed, said apparatus comprising:

10 means for determining, if a second thread attempts to acquire a lock on a specific object that is held by a first thread, whether a bit that represents said lock type on said specific object represents said first lock type;

15 means for setting a contention bit if said bit represents said first lock type;

means for determining, before said first thread unlocks said specific object, whether said bit that represents said lock type represents said first lock type;

20 means for storing in said storage area a special identifier that differs from the identifiers for said plurality of threads;

means for issuing a synchronization command for said storage area;



6. The lock management apparatus according to claim 5, further comprising:

5 means for shifting said first thread, when said contention bit has been set, to an exclusive control state for a mechanism that enables the exclusive control of the accessing of said object, and a thread waiting operation and the transmission to a waiting thread of a notification, both of which are to be performed when a predetermined condition has been established;

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means for permitting said first thread to transmit said notification to said waiting thread;

means for setting said second thread in the busy waiting state, when said predetermined condition has not been established and when said special identifier has been stored, until a thread that maintains said lock on said specific object is no longer present and until said bit that represents said lock type represents said first lock type; and

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20 means for permitting said first thread to exit said exclusive control state.





9. In a shared memory model system, a method whereby, in a state wherein a plurality of threads exist, a bit that represents a lock is stored in a storage area that corresponds to an object, and a queue of a thread that accesses said object is stored to manage a lock on an object, said method comprising:

determining, when a second thread attempts to acquire a lock on a specific object that a first thread has locked, whether a bit that is used to represent said lock on said object represents the locked state;

changing data for the number of queues of threads that access said specific object and storing the updated data when said bit represents said locked state;

storing said second thread in a queue, and shifting said second thread to a control state, for a mechanism that performs a waiting operation for accessing said specific object and a recovery operation by transmitting a notification;

storing said bit that represents said locked state in said storage area before said first thread unlocks said object;

determining whether a thread that is stored in a queue is present;

shifting said first thread to a notification state, wherein said transmission of a notification to said thread that is waiting is initiated, when a thread that is stored in a queue is present; and

5           permitting said first thread to exit said notification state.

10. The lock management method according to claim 9, further comprising:

10           increasing, when said bit that represents said locked state is set, the number of queues of threads that can access said specific object and storing the updated number, and determining whether said bit that represents said lock on said specific object represents said locked state; and

15           reducing, when said bit that represents said locked state is not set, the number of said queues of said threads that access said specific object and storing the updated number, and terminating a locking process without any other process being performed.

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11. In a shared memory model system, a method whereby, in a state wherein a plurality of threads exist, a bit that represents a lock is stored in a storage area that corresponds to an object, and a queue of threads that  
5 access said object is stored to manage a lock on an object, said method comprising:

determining, when a second thread attempts to acquire a lock on a specific object that a first thread has locked, whether a bit that represents said  
10 lock on said object represents the locked state;

changing, when said bit represents said locked state, data for the number of queues of threads that can access said specific object and storing the updated data, and thereafter issuing a  
15 synchronization command for said storage area;

storing said second thread in a queue, and shifting said second thread to a control state for a mechanism that performs a waiting operation, for accessing said specific object, and a recovery  
20 operation by transmitting a notification;

storing in said storage area, before said first thread unlocks said object, said bit that represents said locked state and an identifier that is not related to the representation of said locked state or  
25 an unlocked state;

issuing a synchronization command for said storage area;

storing, in said storage area, data that does not represent said lock on said specific object;

5 determining whether a thread that is stored in a queue is present;

shifting, when a thread that is stored in a queue is present, said first thread to a notification state wherein said transmission is initiated for issuing a notification to said thread that is waiting; and  
10

permitting said first thread to exit said notification state.

12. The lock management method according to claim 11,  
further comprising:

increasing, when said bit that represents said  
locked state is set, the number of queues of threads  
that can access said specific object and storing the  
updated number, and determining whether said bit that  
represents said lock on said specific object  
represents said locked state; and

reducing, when said bit that represents said  
locked state is not set, the number of said queues of  
said threads that access said specific object and  
storing the updated number, and terminating a locking  
process without any other process being performed.

13. The lock management method according to claim 12,  
further comprising:

permitting said second thread, when said bit that  
represents said locked state is set and when an  
identifier that is not related to the representation  
of said locked state or said unlocked state is stored  
in said storage area, to remain in a busy waiting  
state until a thread that maintains said lock on said  
object is no longer present and said bit that  
represents said locked state is changed to represent  
said unlocked state.

14. In a shared memory model system, an apparatus where, in a state wherein a plurality of threads exist, a bit that represents a lock is stored in a storage area that corresponds to an object, and a queue of a thread that accesses said object is stored to manage a lock on an object, said apparatus comprising:

means for determining, when a second thread attempts to acquire a lock on a specific object that a first thread has locked, whether a bit that is used to represent said lock on said object represents the locked state;

means for changing data for the number of queues of threads that access said specific object and storing the updated data when said bit represents said locked state;

means for storing said second thread in a queue, and shifting said second thread to a control state, for a mechanism that performs a waiting operation for accessing said specific object and a recovery operation by transmitting a notification;

means for storing said bit that represents said locked state in said storage area before said first thread unlocks said object;

means for determining whether a thread that is stored in a queue is present;

means for shifting said first thread to a notification state, wherein said transmission of a notification to said thread that is waiting is initiated, when a thread that is stored in a queue is present; and

means for permitting said first thread to exit said notification state.

15. The lock management apparatus according to claim 14, further comprising:

means for increasing, when said bit that represents said locked state is set, the number of queues of threads that can access said specific object and storing the updated number, and determining whether said bit that represents said lock on said specific object represents said locked state; and

means for reducing, when said bit that represents said locked state is not set, the number of said queues of said threads that access said specific object and storing the updated number, and terminating a locking process without any other process being performed.

16. In a shared memory model system, an apparatus where, in a state wherein a plurality of threads exist, a bit that represents a lock is stored in a storage area that corresponds to an object, and a queue of threads that access said object is stored to manage a lock on an object, said apparatus comprising:

means for determining, when a second thread attempts to acquire a lock on a specific object that a first thread has locked, whether a bit that represents said lock on said object represents the locked state;

means for changing, when said bit represents said locked state, data for the number of queues of threads that can access said specific object and storing the updated data, and thereafter issuing a synchronization command for said storage area;

means for storing said second thread in a queue, and shifting said second thread to a control state for a mechanism that performs a waiting operation, for accessing said specific object, and a recovery operation by transmitting a notification;

means for storing in said storage area, before said first thread unlocks said object, said bit that represents said locked state and an identifier that is not related to the representation of said locked state or an unlocked state;



means for issuing a synchronization command for  
said storage area;

5 means for storing, in said storage area, data  
that does not represent said lock on said specific  
object;

means for determining whether a thread that is  
stored in a queue is present;

10 means for shifting, when a thread that is stored  
in a queue is present, said first thread to a  
notification state wherein said transmission is  
initiated for issuing a notification to said thread  
that is waiting; and

15 means for permitting said first thread to exit  
said notification state.

17. The lock management apparatus according to claim 16, further comprising:

means for increasing, when said bit that represents said locked state is set, the number of queues of threads that can access said specific object and storing the updated number, and determining whether said bit that represents said lock on said specific object represents said locked state; and

means for reducing, when said bit that represents said locked state is not set, the number of said queues of said threads that access said specific object and storing the updated number, and terminating a locking process without any other process being performed.

18. The lock management apparatus according to claim 17, further comprising:

5 means for permitting said second thread, when said bit that represents said locked state is set and when an identifier that is not related to the representation of said locked state or said unlocked state is stored in said storage area, to remain in a busy waiting state until a thread that maintains said lock on said object is no longer present and said bit that represents said locked state is changed to  
10 represent said unlocked state.

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